

RANKS OF 0-1 ARRAYS OF SIZE $2 \times 2 \times 2$ AND $2 \times 2 \times 2 \times 2$

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ABSTRACT. We use computer algebra to determine the ranks of arrays of size $2 \times 2 \times 2$ and $2 \times 2 \times 2 \times 2$ with entries in the set $\{0, 1\}$ regarded as a field with two elements, as a Boolean algebra, and as non-negative integers. In the field case we also determine the canonical forms of the arrays with respect to the action of the direct product of the general linear groups.

1. INTRODUCTION

Multidimensional arrays and the related topic of hyperdeterminants have connections with algebraic geometry and representation theory, and applications in numerical analysis, signal processing, chemometrics and psychometrics. For the connections with algebraic geometry see [5, 6, 10]; for a connection with representation theory see [1]. For surveys of the applications, see [2, 8, 9, 11].

Most research on this topic assumes that the arrays have entries in \mathbb{R} or \mathbb{C} , the fields of real and complex numbers. In this paper, we consider arrays with entries in $\{0, 1\}$, which can be regarded as the field \mathbb{F}_2 with two elements ($1 + 1 = 0$), as a Boolean algebra ($1 + 1 = 1$), or as non-negative integers ($1 + 1 = 2$). In this case, the total number of arrays is finite, and the problem of determining the rank of an array reduces to combinatorial enumeration which can be performed by computer. We obtain a classification by rank of all $2 \times 2 \times 2$ and $2 \times 2 \times 2 \times 2$ arrays in these three cases. In the field case we determine the canonical forms of the arrays with respect to the action of the finite groups $GL_2(\mathbb{F}_2)^3$ and $GL_2(\mathbb{F}_2)^4$ and the extended groups $GL_2(\mathbb{F}_2)^3 \rtimes S_3$ and $GL_2(\mathbb{F}_2)^4 \rtimes S_4$.

The canonical forms of $2 \times 2 \times 2$ arrays have been determined over the real numbers [3] and over the complex numbers [4, 5]. Analogous results for $2 \times 2 \times 2 \times 2$ arrays over \mathbb{R} or \mathbb{C} have not yet been found, but see [7]. Since the results for $2 \times 2 \times 2$ arrays over \mathbb{R} and \mathbb{C} are similar to the results in this paper over \mathbb{F}_2 , we hope that our results for $2 \times 2 \times 2 \times 2$ arrays over \mathbb{F}_2 will provide some useful information towards a classification of canonical forms of $2 \times 2 \times 2 \times 2$ arrays over \mathbb{R} and \mathbb{C} .

2. PRELIMINARIES

We consider n -dimensional arrays of size $2 \times \cdots \times 2$ (n factors, $n = 3, 4$) with entries in $\{0, 1\}$.

Definition 1. The **flattening** of the array $X = (x_{i_1 \dots i_n})$, $1 \leq i_1, \dots, i_n \leq 2$, is

$$\text{flat}(X) = \begin{bmatrix} x_{1\dots 1} & \cdots & x_{i_1 \dots i_n} & \cdots & x_{2\dots 2} \end{bmatrix},$$

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where the entries are in lexicographical order of the n -tuples of subscripts: $i_1 \cdots i_n$ precedes $i'_1 \cdots i'_n$ if and only if $i_j < i'_j$ where j is the least index with $i_j \neq i'_j$.

Definition 2. If X and Y are two arrays then X **precedes** Y if $\text{flat}(X)$ precedes $\text{flat}(Y)$ in lexicographical order: that is, $x_{i_1 \cdots i_n} < y_{i_1 \cdots i_n}$ where $i_1 \cdots i_n$ is the least n -tuple with $x_{i_1 \cdots i_n} \neq y_{i_1 \cdots i_n}$. The **minimal element** of a set of arrays is defined with respect to this total order.

Definition 3. There are three nonzero 2-dimensional vectors: $[0, 1]$, $[1, 0]$, $[1, 1]$. The **outer product** $X = V_1 \otimes \cdots \otimes V_n$ of n vectors $V_j = [v_{j1}, v_{j2}]$, $1 \leq j \leq n$, is

$$X = (x_{i_1 \cdots i_n}), \quad x_{i_1 \cdots i_n} = v_{1i_1} \cdots v_{ni_n}.$$

Definition 4. The **rank** of $X = (x_{i_1 \cdots i_n})$ is the minimal number R of terms in the expression of X as a sum of outer products:

$$X = \sum_{r=1}^R V_1^{(r)} \otimes \cdots \otimes V_n^{(r)}.$$

The definition of addition depends on the algebraic structure of $\{0, 1\}$. For the field with two elements, $1 + 1 = 0$; for the Boolean algebra, $1 + 1 = 1$; for non-negative integers, $1 + 1 = 2$, and this means that we exclude any sums in which two outer products both have an entry 1 in the same position.

Lemma 5. *The only array of rank 0 is the zero array. An array of rank 1 is the same as an outer product of nonzero vectors, and there are 3^n such arrays. The total number of n -dimensional $2 \times \cdots \times 2$ arrays with entries in $\{0, 1\}$ is 2^{2^n} .*

Algorithm 6. Fix a dimension n . Assume that we have already computed the arrays of rank r . To compute the arrays of rank $r + 1$, we consider all sums $X + Y$ where $\text{rank}(X) = r$ and $\text{rank}(Y) = 1$. Clearly $\text{rank}(X + Y) \leq r + 1$, but it is possible that $\text{rank}(X + Y) \leq r$, so we only retain those $X + Y$ which have not already been computed: the arrays which have rank exactly $r + 1$. This algorithm is presented in pseudocode for $n = 3$ in Table 1.

If $\{0, 1\}$ is the field \mathbb{F}_2 with two elements then we consider the action of the direct product of n copies of the general linear group $GL_2(\mathbb{F}_2)$.

Definition 7. The group $GL_2(\mathbb{F}_2)$ consists of six matrices in lexicographical order:

$$\begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix}, \quad \begin{bmatrix} 0 & 1 \\ 1 & 1 \end{bmatrix}, \quad \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}, \quad \begin{bmatrix} 1 & 0 \\ 1 & 1 \end{bmatrix}, \quad \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}, \quad \begin{bmatrix} 1 & 1 \\ 1 & 0 \end{bmatrix}.$$

Lemma 8. *The group $GL_2(\mathbb{F}_2)$ is isomorphic to S_3 permuting the nonzero vectors.*

Definition 9. The **small symmetry group** of $2 \times \cdots \times 2$ arrays over \mathbb{F}_2 is the direct product $GL_2(\mathbb{F}_2)^n$ acting by simultaneous changes of basis along the n directions. The **large symmetry group** of these arrays is the semi-direct product $GL_2(\mathbb{F}_2)^n \rtimes S_n$ where S_n acts by permuting the n copies of $GL_2(\mathbb{F}_2)$. The element $A \in GL_2(\mathbb{F}_2)$ acts along the first direction of an array $X = (x_{i_1 \cdots i_n})$ as follows: for each $(n-1)$ -tuple $i_2 \cdots i_n$ we consider the column vector

$$V_{i_2 \cdots i_n} = [x_{1i_2 \cdots i_n}, x_{2i_2 \cdots i_n}]^t \in \mathbb{F}_2^2,$$

and compute $AV_{i_2 \cdots i_n} = [y_{1i_2 \cdots i_n}, y_{2i_2 \cdots i_n}]^t \in \mathbb{F}_2^2$; then we define the array $A \cdot X$ to be $Y = (y_{i_1 \cdots i_n})$. The actions along the other $n - 1$ directions are similar.

```

flatten( $x$ )
  return( $[x_{111}, x_{112}, x_{121}, x_{122}, x_{211}, x_{212}, x_{221}, x_{222}]$ )

outerproduct( $a, b, c$ )
  for  $i = 1, 2$  do for  $j = 1, 2$  do for  $k = 1, 2$  do:
    set  $x_{ijk} \leftarrow a_i b_j c_k$ 
  return( $x$ )

• set  $\text{vectors} \leftarrow \{[1, 0], [0, 1], [1, 1]\}$ 
• set  $\text{arrayset}[0] \leftarrow \{[0, 0, 0, 0, 0, 0, 0, 0]\}$ 
• set  $\text{arrayset}[1] \leftarrow \{\}$ 
• for  $a$  in  $\text{vectors}$  do for  $b$  in  $\text{vectors}$  do for  $c$  in  $\text{vectors}$  do
  – set  $x \leftarrow \text{flatten}(\text{outerproduct}(a, b, c))$ 
  – if  $x \notin \text{arrayset}[0]$  and  $x \notin \text{arrayset}[1]$  then
    set  $\text{arrayset}[1] \leftarrow \text{arrayset}[1] \cup \{x\}$ 
• set  $r \leftarrow 1$ 
• while  $\text{arrayset}[r] \neq \{\}$  do:
  – set  $\text{arrayset}[r+1] \leftarrow \{\}$ 
  – for  $x \in \text{arrayset}[r]$  do for  $y \in \text{arrayset}[1]$  do
    * set  $z \leftarrow [x_1+y_1, \dots, x_8+y_8]$ 
    * if  $z \notin \text{arrayset}[s]$  for  $s = 0, \dots, r+1$  then
      set  $\text{arrayset}[r+1] \leftarrow \text{arrayset}[r+1] \cup \{z\}$ 
  – set  $r \leftarrow r + 1$ 
• set  $\text{maximumrank} \leftarrow r - 1$ 

```

TABLE 1. Algorithm 6 in pseudocode

Lemma 10. *The actions of the symmetry groups do not change the rank.*

Proof. de Silva and Lim [3, Lemma 2.3, page 1092] applies to any field. □

The actions of the symmetry groups decompose the set of n -dimensional arrays into a disjoint union of orbits; the arrays in each orbit are equivalent under the group action.

Algorithm 11. Fix a dimension n and consider $2 \times \dots \times 2$ arrays over \mathbb{F}_2 . Assume we have computed the set of arrays for each possible rank, and that these sets are totally ordered. For each rank, we perform the following iteration:

- Choose the minimal element of the set of arrays.
- Compute the orbit of this element under the action of the symmetry group.
- Remove the elements of this orbit from the set of arrays of the given rank.

This iteration terminates when there are no more arrays of the given rank. This algorithm is presented in pseudocode for $n = 3$ in Table 2.

Definition 12. The minimal element in each orbit is the **canonical form** of the arrays in that orbit.

```

unflatten( $x$ )
  set  $t \leftarrow 0$ 
  for  $i = 1, 2$  do for  $j = 1, 2$  do for  $k = 1, 2$  do:
    set  $t \leftarrow t + 1$ 
    set  $y_{ijk} \leftarrow x_t$ 
  return( $y$ )

groupaction( $g, x, m$ )
  set  $y \leftarrow \text{unflatten}(x)$ 
  if  $m = 1$  then
    for  $j = 1, 2$  do for  $k = 1, 2$  do
      set  $v \leftarrow [y_{1jk}, y_{2jk}]$ 
      set  $w \leftarrow [g_{11}v_1 + g_{12}v_2, g_{21}v_1 + g_{22}v_2]$ 
      for  $i = 1, 2$  do: set  $y_{ijk} \leftarrow w_i$ 
  if  $m = 2$  then ... (similar for second subscript)
  if  $m = 3$  then ... (similar for third subscript)
  return(flatten( $y$ ))

smallorbit( $x$ )
  set result  $\leftarrow \{\}$ 
  for  $a \in GL_2(\mathbb{F}_2)$  do:
    set  $y \leftarrow \text{groupaction}(a, x, 1)$ 
    for  $b \in GL_2(\mathbb{F}_2)$  do:
      set  $z \leftarrow \text{groupaction}(b, y, 2)$ 
      for  $c \in GL_2(\mathbb{F}_2)$  do:
        set  $w \leftarrow \text{groupaction}(c, z, 3)$ 
        set result  $\leftarrow \text{result} \cup \{w\}$ 
  return(result)

largeorbit( $x$ )
  set  $y \leftarrow \text{unflatten}(x)$ 
  set result  $\leftarrow \{\}$ 
  for  $p \in S_3$  do:
    for  $i = 1, 2$  do for  $j = 1, 2$  do for  $k = 1, 2$  do:
      set  $m \leftarrow [i, j, k]$ 
      set  $z_{ijk} \leftarrow y_{m_{p(1)}m_{p(2)}m_{p(3)}}$ 
      set result  $\leftarrow \text{result} \cup \text{smallorbit}(\text{flatten}(z))$ 
  return(result)

• for  $r = 0, \dots, \text{maximumrank}$  do:
  set representatives[ $r$ ]  $\leftarrow \{\}$ 
  set remaining  $\leftarrow \text{arrayset}[r]$ 
  while remaining  $\neq \{\}$  do:
    set  $x \leftarrow \text{remaining}[1]$ 
    set xorbit  $\leftarrow \text{largeorbit}(x)$ 
    append xorbit[1] to representatives[ $r$ ]
    set remaining  $\leftarrow \text{remaining} \setminus \text{xorbit}$ 

```

TABLE 2. Algorithm 11 in pseudocode

rank	orbit size	canonical form
0	1	$\left[\begin{array}{cc cc} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{array} \right]$
1	27	$\left[\begin{array}{cc cc} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{array} \right]$
2	54	$\left[\begin{array}{cc cc} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{array} \right]$
2	108	$\left[\begin{array}{cc cc} 0 & 0 & 1 & 0 \\ 0 & 1 & 0 & 0 \end{array} \right]$
3	54	$\left[\begin{array}{cc cc} 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{array} \right]$
3	12	$\left[\begin{array}{cc cc} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{array} \right]$

TABLE 3. Large orbits of $2 \times 2 \times 2$ arrays over \mathbb{F}_2

Lemma 13. *Lower bounds for the number of canonical forms for the small symmetry group and the large symmetry group are respectively*

$$\left\lceil \frac{2^{2^n}}{6^n} \right\rceil, \quad \left\lceil \frac{2^{2^n}}{6^n n!} \right\rceil.$$

Proof. Lemma 5 shows that there are 2^{2^n} such arrays. Definition 9 implies that the small and large symmetry groups have orders 6^n and $6^n n!$ respectively. The claim follows from the theory of group actions on a finite set. \square

Remark 14. For $3 \leq n \leq 6$, we have the following lower bounds for the number of orbits for the small and large symmetry groups. From this it is clear that complete results will not be publishable for $n \geq 5$:

n	lower bound (small group)	lower bound (large group)
3	2	1
4	51	3
5	552337	4603
6	395377745064077	549135757034

3. ARRAYS OF SIZE $2 \times 2 \times 2$

The set of $2 \times 2 \times 2$ arrays with entries in $\{0, 1\}$ contains 256 elements. We represent such an array $X = (x_{ijk})$ in the matrix form

$$\text{Mat}(X) = \left[\begin{array}{cc|cc} x_{111} & x_{121} & x_{112} & x_{122} \\ x_{211} & x_{221} & x_{212} & x_{222} \end{array} \right],$$

where the third subscript distinguishes the left and right blocks, which are the first and second frontal slices.

rank	ones	number	representative
0	0	1	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 0 \end{bmatrix}$
1	1	8	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix}$
1	2	12	$\begin{bmatrix} 0 & 0 & 0 & 0 \\ 0 & 0 & 1 & 1 \end{bmatrix}$
1	4	6	$\begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$
1	8	1	$\begin{bmatrix} 1 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$
2	2	16	$\begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{bmatrix}$
2	3	48	$\begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 \end{bmatrix}$
2	4	30	$\begin{bmatrix} 0 & 0 & 1 & 0 \\ 0 & 1 & 1 & 1 \end{bmatrix}$
2	5	24	$\begin{bmatrix} 0 & 0 & 1 & 1 \\ 0 & 1 & 1 & 1 \end{bmatrix}$
2	6	12	$\begin{bmatrix} 0 & 0 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$
3	3	8	$\begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 0 \end{bmatrix}$
3	4	32	$\begin{bmatrix} 0 & 0 & 0 & 1 \\ 0 & 1 & 1 & 1 \end{bmatrix}$
3	5	24	$\begin{bmatrix} 0 & 0 & 1 & 1 \\ 1 & 1 & 0 & 1 \end{bmatrix}$
3	6	16	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 0 & 1 & 1 \end{bmatrix}$
3	7	8	$\begin{bmatrix} 0 & 1 & 1 & 1 \\ 1 & 1 & 1 & 1 \end{bmatrix}$
4	4	2	$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 0 & 1 \end{bmatrix}$
4	5	8	$\begin{bmatrix} 0 & 1 & 1 & 0 \\ 1 & 0 & 1 & 1 \end{bmatrix}$

TABLE 4. Ranks and minimal representatives for $2 \times 2 \times 2$ Boolean arrays

3.1. The field with two elements ($1 + 1 = 0$). Algorithm 6 shows that in this case the maximum rank is 3; the number of arrays of each rank is 1, 27, 162, 66. The percentages of ranks 0, 1, 2, 3 are approximately 0, 11, 63, 26; in contrast, the percentages over \mathbb{R} are approximately 0, 0, 79, 21. For the large symmetry group

$GL_2(\mathbb{F}_2)^3 \rtimes S_3$, the ranks, orbit sizes, and canonical forms are given in Table 3. For the small symmetry group $GL_2(\mathbb{F}_2)^3$, the first orbit in rank 2 splits into three orbits each of size 18 with canonical forms

$$\left[\begin{array}{cc|cc} 0 & 0 & 0 & 1 \\ 0 & 0 & 1 & 0 \end{array} \right], \quad \left[\begin{array}{cc|cc} 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 0 \end{array} \right], \quad \left[\begin{array}{cc|cc} 0 & 0 & 0 & 1 \\ 0 & 1 & 0 & 0 \end{array} \right].$$

For the small symmetry group, there are eight orbits, the same as in the real case.

3.2. The Boolean algebra ($1 + 1 = 1$). The maximum rank is 4; the number of arrays of each rank is 1, 27, 130, 88, 10. The percentages of ranks 0, 1, 2, 3, 4 are approximately 0, 11, 51, 34, 4. Instead of canonical forms for a group action, which do not exist in the Boolean case, we partition the arrays in each rank by the number of entries equal to 1; the results are given in Table 4.

3.3. Non-negative integers ($1 + 1 = 2$). The results are the same as in the Boolean case. This has the corollary that every $2 \times 2 \times 2$ Boolean array of rank r can be written as the sum of r outer products such that no two terms have an entry 1 in the same position. (When we consider arrays of size $2 \times 2 \times 2 \times 2$, the Boolean and integer cases will no longer be identical.)

4. ARRAYS OF SIZE $2 \times 2 \times 2 \times 2$

The set of $2 \times 2 \times 2 \times 2$ arrays with entries in $\{0, 1\}$ contains 65536 elements.

4.1. The field with two elements ($1 + 1 = 0$). Algorithm 6 shows that in this case the maximum rank is 6. The number of arrays of each rank and the approximate percentages are as follows:

rank	0	1	2	3	4	5	6
number	1	81	2268	21744	37530	3888	24
$\approx \%$	0.002	0.124	3.461	33.179	57.266	5.933	0.037

For the large symmetry group $GL_2(\mathbb{F}_2)^4 \rtimes S_4$, there are 30 orbits; the ranks, orbit sizes, and canonical forms are given in Table 5. For the small symmetry group $GL_2(\mathbb{F}_2)^4$, there are 112 orbits. The large orbits split into small orbits as follows, where we mention only those large orbits that are not small orbits, and write $x \rightarrow y \cdot z$ to indicate that large orbit x splits into y small orbits each of size z :

rank 2	rank 3	rank 4	rank 5
3 \rightarrow 6 \cdot 54	6 \rightarrow 4 \cdot 162	15 \rightarrow 4 \cdot 648	26 \rightarrow 6 \cdot 108
4 \rightarrow 4 \cdot 324	7 \rightarrow 4 \cdot 36	16 \rightarrow 4 \cdot 1296	
	8 \rightarrow 6 \cdot 648	17 \rightarrow 3 \cdot 36	
	9 \rightarrow 4 \cdot 648	18 \rightarrow 3 \cdot 324	
	10 \rightarrow 4 \cdot 648	19 \rightarrow 6 \cdot 324	
	11 \rightarrow 3 \cdot 1296	20 \rightarrow 3 \cdot 648	
	12 \rightarrow 6 \cdot 1296	21 \rightarrow 12 \cdot 648	
		22 \rightarrow 6 \cdot 216	
		23 \rightarrow 6 \cdot 1296	
		24 \rightarrow 3 \cdot 1296	
		25 \rightarrow 6 \cdot 648	

4.2. The Boolean algebra ($1 + 1 = 1$). The maximum rank is 8. The number of arrays of each rank and the approximate percentages are as follows:

rank	0	1	2	3	4	5	6	7	8
number	1	81	1804	13472	28904	17032	3704	512	26
$\approx \%$	0.002	0.124	2.753	20.557	44.104	25.989	5.652	0.781	0.04

The results are given in Table 6 by rank and number of entries equal to 1.

4.3. Non-negative integers ($1 + 1 = 2$). The maximum rank is 8. The number of arrays of each rank and the approximate percentages are as follows:

rank	0	1	2	3	4	5	6	7	8
number	1	81	1756	12848	28788	17568	3908	560	26
$\approx \%$	0.002	0.124	2.679	19.604	43.927	26.807	5.963	0.854	0.04

The results are given in Table 7 by rank and number of entries equal to 1.

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	rank	large orbit size	canonical form (flattened)
1	0	1	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0
2	1	81	0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 0 1
3	2	324	0 0 0 0 0 0 0 0 0 0 0 0 0 0 1 1 0
4	2	1296	0 0 0 0 0 0 0 0 0 0 0 0 1 1 0 0 0
5	2	648	0 0 0 0 0 0 0 0 1 1 0 0 0 0 0 0 0
6	3	648	0 0 0 0 0 0 0 0 0 0 0 0 1 0 1 1 0
7	3	144	0 0 0 0 0 0 0 0 0 1 1 0 1 0 1 1 1
8	3	3888	0 0 0 0 0 0 0 0 1 0 0 0 1 1 0 0 0
9	3	2592	0 0 0 0 0 0 0 0 1 0 0 1 0 1 1 0 0
10	3	2592	0 0 0 0 0 0 0 0 1 0 1 1 0 1 0 1 0
11	3	3888	0 0 0 0 0 0 0 0 1 1 0 0 0 0 0 1 0
12	3	7776	0 0 0 0 0 0 0 0 1 1 0 0 0 0 1 1 0
13	3	216	0 0 0 1 1 0 0 0 1 1 1 0 1 1 1 1 1
14	4	162	0 0 0 0 0 0 0 0 1 0 0 0 1 0 1 1 0
15	4	2592	0 0 0 0 0 0 0 0 1 0 1 1 0 1 0 0 0
16	4	5184	0 0 0 0 0 0 0 0 1 1 0 0 1 0 1 1 0
17	4	108	0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 0
18	4	972	0 0 0 0 0 1 1 0 0 1 1 0 0 0 0 0 1
19	4	1944	0 0 0 0 0 1 1 0 0 1 1 0 0 0 1 0 0
20	4	1944	0 0 0 0 0 1 1 0 0 1 1 1 0 0 1 0 0
21	4	7776	0 0 0 0 0 1 1 0 0 1 1 1 1 0 0 0 0
22	4	1296	0 0 0 0 0 1 1 0 1 0 1 1 0 0 0 0 0
23	4	7776	0 0 0 0 0 1 1 0 1 0 1 1 0 0 0 0 1
24	4	3888	0 0 0 1 0 1 1 0 1 0 0 0 0 0 0 1 1
25	4	3888	0 0 0 1 0 1 1 0 1 0 0 0 1 0 1 1 1
26	5	648	0 0 0 0 0 1 1 0 0 1 1 0 1 0 1 1 1
27	5	648	0 0 0 1 0 1 1 0 0 1 1 0 1 0 0 0 0
28	5	1296	0 0 0 1 0 1 1 0 0 1 1 0 1 0 0 1 0
29	5	1296	0 0 0 1 0 1 1 0 1 0 0 0 0 0 0 0 1
30	6	24	0 1 1 0 1 0 1 1 1 0 1 1 1 1 0 1 1

TABLE 5. Large orbits of $2 \times 2 \times 2 \times 2$ arrays over \mathbb{F}_2

	rank	ones	size	representative
1	0	0	1	0000000000000000
2	1	1	16	0000000000000001
3	1	2	32	0000000000000011
4	1	4	24	0000000000000111
5	1	8	8	0000000011111111
6	1	16	1	1111111111111111
7	2	2	88	0000000000000110
8	2	3	352	0000000000000111
9	2	4	352	00000000000011011
10	2	5	288	00000000000011111
11	2	6	384	00000000001111111
12	2	7	48	00000011010101111
13	2	8	108	00000011110011111
14	2	9	64	00000001111111111
15	2	10	96	00000011111111111
16	2	12	24	00001111111111111
17	3	3	208	0000000000010110
18	3	4	1216	0000000000010111
19	3	5	2304	00000000000111101
20	3	6	2512	00000000001101111
21	3	7	2656	00000000001111111
22	3	8	1904	00000000111101111
23	3	9	1056	00000011110111111
24	3	10	656	00000110111111111
25	3	11	576	00000111111111111
26	3	12	256	00011011111111111
27	3	13	96	00011111111111111
28	3	14	32	00111111111111111
29	4	4	228	0000000001101001
30	4	5	1648	0000000001101011
31	4	6	4048	00000000100111110
32	4	7	5856	00000000101101111
33	4	8	6304	00000000101111111
34	4	9	5200	00000001101111111
35	4	10	3200	00000011110111111
36	4	11	1408	00001111111101111
37	4	12	652	00010111111111111
38	4	13	256	00111101111111111
39	4	14	88	01101111111111111
40	4	15	16	01111111111111111
41	5	5	128	00000000110010110
42	5	6	1008	00000000110010111
43	5	7	2416	00000001101101101
44	5	8	3568	00000011001101111
45	5	9	4016	00000011001111111
46	5	10	3088	00000011101111111
47	5	11	1888	00010111111101111
48	5	12	712	00011111111101111
49	5	13	208	01101011111111111
50	6	6	56	00000011001101001
51	6	7	448	00000011001101011
52	6	8	848	00000011101111001
53	6	9	928	00010110011011111
54	6	10	848	00010110011111111
55	6	11	416	00010111011111111
56	6	12	160	01101011110111111
57	7	7	16	0001011001101001
58	7	8	128	0001011001101011
59	7	9	160	0001011111101001
60	7	10	112	0011110111010110
61	7	11	80	01101001101111111
62	7	12	16	01101011101111111
63	8	8	2	0110100110010110
64	8	9	16	0110100110010111
65	8	10	8	0110101111010110

TABLE 6. Minimal representatives for $2 \times 2 \times 2 \times 2$ Boolean arrays

	rank	ones	size	representative
1	0	0	1	0000000000000000
2	1	1	16	0000000000000001
3	1	2	32	0000000000000011
4	1	4	24	0000000000000111
5	1	8	8	0000000011111111
6	1	16	1	1111111111111111
7	2	2	88	0000000000000110
8	2	3	352	0000000000000111
9	2	4	352	00000000000011011
10	2	5	288	00000000000011111
11	2	6	384	00000000001111111
12	2	8	108	00000011110011111
13	2	9	64	00000001111111111
14	2	10	96	00000011111111111
15	2	12	24	00001111111111111
16	3	3	208	00000000000101110
17	3	4	1216	00000000000101111
18	3	5	2304	00000000000111101
19	3	6	2512	00000000001101111
20	3	7	2704	00000000001111111
21	3	8	1664	00000001111011111
22	3	9	864	00000011110111111
23	3	10	608	00000110111111111
24	3	11	384	00000111111111111
25	3	12	256	00011011111111111
26	3	13	96	00011111111111111
27	3	14	32	00111111111111111
28	4	4	228	00000000011010011
29	4	5	1648	00000000011010111
30	4	6	4048	00000001001111110
31	4	7	5856	00000001011011111
32	4	8	6544	00000001011111111
33	4	9	5104	00000011011111111
34	4	10	3056	00000111101111111
35	4	11	1504	00001111111101111
36	4	12	448	00010111111111111
37	4	13	256	00111101111111111
38	4	14	80	01101111111111111
39	4	15	16	01111111111111111
40	5	5	128	00000001100101110
41	5	6	1008	00000001100101111
42	5	7	2416	00000011011011011
43	5	8	3568	00000110011011111
44	5	9	4304	00000110011111111
45	5	10	3088	00000111011111111
46	5	11	1984	00010111111101111
47	5	12	904	00011111111101111
48	5	13	160	01101011111111111
49	5	14	8	01111111111111110
50	6	6	56	00000110011010011
51	6	7	448	00000110011010111
52	6	8	848	00000111011110011
53	6	9	928	00010110011011111
54	6	10	1040	00010110011111111
55	6	11	368	00010111011111111
56	6	12	172	01101011110111111
57	6	13	48	01101111111101111
58	7	7	16	00010110011010011
59	7	8	128	00010110011010111
60	7	9	160	00010111111010011
61	7	10	112	00111101110101110
62	7	11	128	01101001101111111
63	7	12	16	01101011101111111
64	8	8	2	01101001100101110
65	8	9	16	01101001100101111
66	8	10	8	01101011110101110

TABLE 7. Minimal representatives for $2 \times 2 \times 2 \times 2$ integer arrays